

# Creative Telescoping

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Geometric control and related fields



# What is Creative Telescoping?

Creative telescoping is a methodology

- ▶ to deal with parametrized symbolic sums and integrals
- ▶ that yields differential/recurrence equations for such expressions
- ▶ that became popular in computer algebra in the past 25 years

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$$f_n := \sum_{k=1}^{\infty} \frac{1}{k(k+n)} \rightsquigarrow (n+2)^2 f_{n+2} = (n+1)(2n+3) f_{n+1} - n(n+1) f_n$$

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Creative telescoping can be performed in different settings:

- ▶ difference fields (Carsten Schneider, RISC)
- ▶ differential fields (Clemens Raab, DESY → RICAM)
- ▶ holonomic functions (here)

## Motivating Example

T. Combet, C. K.: Third order integrability conditions for homogeneous potentials of degree  $-1$ . Define:

$$P_n(t) = \frac{1}{t^2 - 1} \frac{\partial^{n-1}}{\partial t^{n-1}} (t^2 - 1)^n$$
$$Q_n(t) = \epsilon_n P_n(t) \operatorname{arctanh} \left( \frac{1}{t} \right) + \frac{W_n(t)}{t^2 - 1}$$

where  $W_n$  are polynomials and  $\epsilon_n$  is a real sequence, given by

$$W_{2k}(t) = \frac{(-1)^k (t^2 - 1)}{2^{4k}} \left( \frac{\pi {}_2F_1\left(\frac{1}{2} - k, k + 1, \frac{1}{2}, t^2\right)}{\Gamma\left(k + \frac{1}{2}\right)^2} + \frac{2kt(2k + 1) \operatorname{arctanh}(t) {}_2F_1\left(1 - k, k + \frac{3}{2}, \frac{3}{2}, t^2\right)}{(k!)^2} \right),$$

$$W_{2k+1}(t) = \dots$$

$$\epsilon_n = \frac{n(n+1)}{4^n (n!)^2}.$$

## Motivating Example

**Task:** Compute recurrences for the following quantities:

$$f_1(n) = \langle \alpha^3 \rangle 2\epsilon_n^{-2} \operatorname{Res}_{t=\infty} \left( (t^2 - 1)^2 (Q_n + \epsilon_n \alpha P_n)^2 (Q_2 + \epsilon_2 \alpha P_2) \right. \\ \left. \times \int (t^2 - 1)^2 (Q_n + \epsilon_n \alpha P_n)^2 P_2 dt \right),$$

$$f_2(n) = \langle \alpha^3 \rangle 2\epsilon_n^{-2} \operatorname{Res}_{t=\infty} \left( (t^2 - 1)^2 (Q_n + \epsilon_n \alpha P_n)^3 \right. \\ \left. \times \int (t^2 - 1)^2 (Q_n + \epsilon_n \alpha P_n)^2 P_n dt \right),$$

$$f_3(n) = \langle \alpha^3 \rangle \frac{1}{6} \epsilon_n^{-2} \operatorname{Res}_{t=\infty} \left( (t^2 - 1)^3 (Q_n + \epsilon_n \alpha P_n)^4 \right).$$

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Method for doing integrals and sums  
(aka Feynman's differentiating under the integral sign)

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Then  $F(n) = \sum_{k=a}^b (g(n, k + 1) - g(n, k)) = g(n, b + 1) - g(n, a)$ .

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**Creative Telescoping:** write

$$c_d(n)f(n + d, k) + \cdots + c_0(n)f(n, k) = g(n, k + 1) - g(n, k).$$

Summing from  $a$  to  $b$  yields a recurrence for  $F(n)$ :

$$c_d(n)F(n + d) + \cdots + c_0(n)F(n) = g(n, b + 1) - g(n, a).$$

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Consider the following integration problem:  $F(x) := \int_a^b f(x, y) dy$

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**Creative Telescoping:** write

$$c_d(x) \frac{d^d}{dx^d} f(x, y) + \cdots + c_0(x) f(x, y) = \frac{d}{dy}g(x, y).$$

Integrating from  $a$  to  $b$  yields a differential equation for  $F(x)$ :

$$c_d(x) \frac{d^d}{dx^d} F(x) + \cdots + c_0(x) F(x) = g(x, b) - g(x, a)$$

## D-finite and P-recursive

**Definition:** A function  $f(x)$  is called **D-finite** if it satisfies a linear ordinary differential equation with polynomial coefficients:

$$p_d(x)f^{(d)}(x) + \cdots + p_1(x)f'(x) + p_0(x)f(x) = 0,$$

$p_0, \dots, p_d \in \mathbb{K}[x]$  not all zero.

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→ In both cases, only finitely many initial conditions!

→ Also called **holonomic function** resp. **holonomic sequence**.

## Example: Harmonic Numbers

**Example:** The harmonic numbers  $H_n = \sum_{k=1}^n \frac{1}{k}$  satisfy the recurrence

$$nH_n = (2n - 1)H_{n-1} - (n - 1)H_{n-2} \quad (n \geq 2)$$

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## Closure Properties

If  $f(x)$  and  $g(x)$  are D-finite then also the following are D-finite

- ▶  $f(x) + g(x)$
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A sequence is P-recursive iff its generating function is D-finite.

## Proof

Show that for P-recursive sequences  $f(n)$  and  $g(n)$  also  $h(n) = f(n)g(n)$  is P-recursive. Assume  $f$  and  $g$  satisfy recurrences of order  $d_1$  and  $d_2$ , respectively.

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Ansatz: want to find  $c_0, \dots, c_d \in \mathbb{K}[n]$  such that

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All coefficients  $r_{i,j}$  must vanish: this yields  $d_1d_2$  equations for the unknowns  $c_0, \dots, c_d$ . The choice  $d = d_1d_2$  ensures a solution.

## Multivariate Generalization

Try to generalize the finiteness property to

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(the  $n_i$  are called **discrete variables**)
- ▶ mixed setting: functions in several continuous and discrete variables  $f(x_1, \dots, x_s, n_1, \dots, n_r)$

## Example: Legendre Polynomials $P_n(x)$

This family of (orthogonal) polynomials is a particular solution of the differential equation

$$(x^2 - 1)P_n''(x) + 2xP_n'(x) - n(n + 1)P_n(x) = 0.$$

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$$P_n^{(4)}(x) = -\frac{6x}{x^2-1}P_n^{(3)}(x) + \frac{(n-2)(n+3)}{x^2-1}P_n''(x)$$



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$$(x^2 - 1)P_n^{(3)}(x) + 4xP_n''(x) - (n - 1)(n + 2)P_n'(x) = 0$$

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This family of (orthogonal) polynomials is a particular solution of the differential equation

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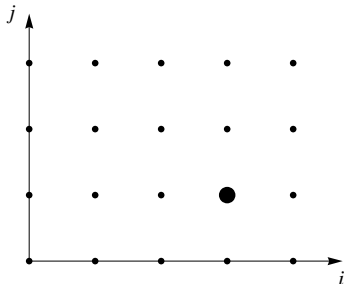
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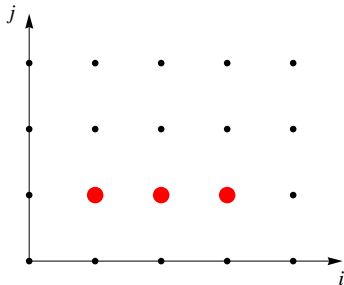
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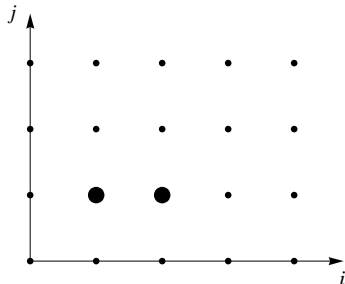
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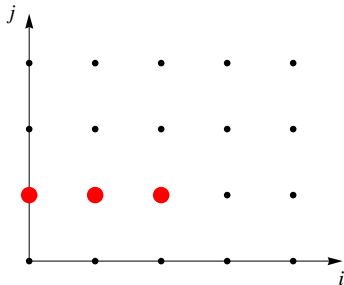
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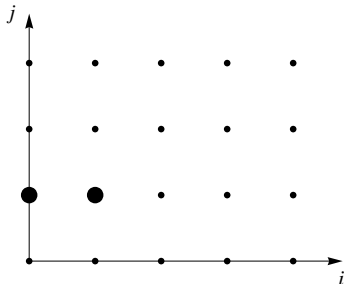
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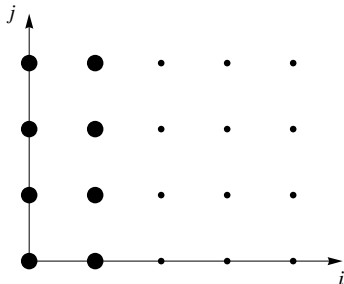
$$P_{n+1}^{(3)}(x) = \frac{(n^2x^2 - n^2 + 3nx^2 - 3n + 8x^2)}{(x^2 - 1)^2} P_{n+1}'(x) - \frac{4(n^2x + 3nx + 2x)}{(x^2 - 1)^2} P_{n+1}(x)$$

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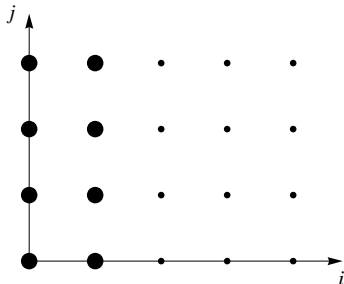
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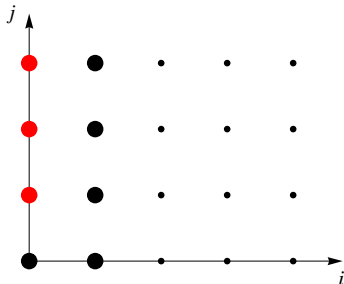
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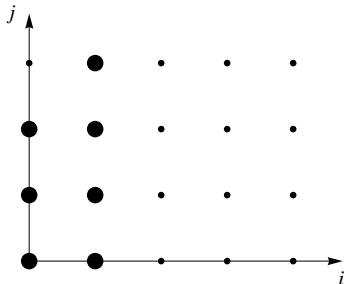
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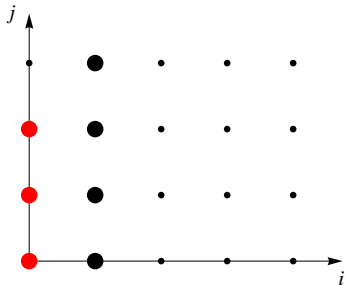
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$$P_{n+3}(x) = \frac{4n^2x^2 - n^2 + 16nx^2 - 4n + 15x^2 - 4}{(n+2)(n+3)} P_{n+1}(x) - \frac{2n^2x + 7nx + 5x}{(n+2)(n+3)} P_n(x)$$

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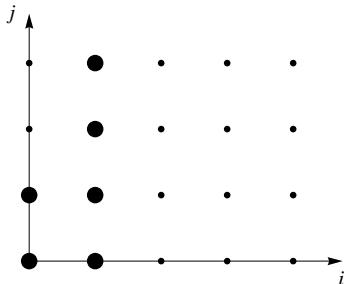
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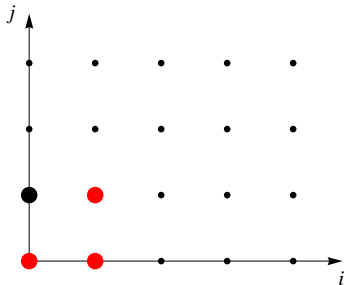
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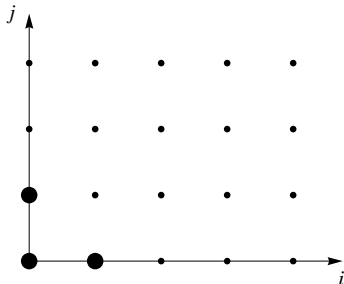
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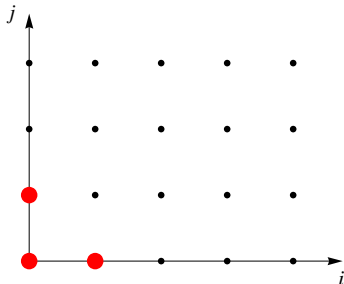
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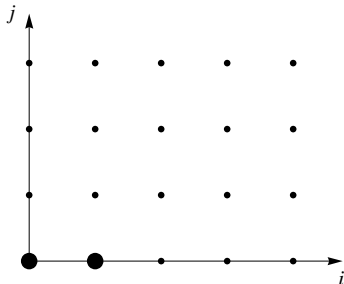
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→  $P_n(x)$  is  $\partial$ -finite w.r.t.  $n$  and  $x$  (of rank 2).

## $\partial$ -Finiteness

First attempt at a definition:

Let  $f(x_1, \dots, x_s, n_1, \dots, n_r)$  be a function in the continuous variables  $x_1, \dots, x_s$  and in the discrete variables  $n_1, \dots, n_r$ .

If there is a finite set of basis functions of the form

$$\frac{d^{i_1}}{dx_1^{i_1}} \cdots \frac{d^{i_s}}{dx_s^{i_s}} f(x_1, \dots, x_s, n_1 + j_1, \dots, n_r + j_r)$$

with  $i_1, \dots, i_s, j_1, \dots, j_r \in \mathbb{N}$  such that any shifted partial derivative of  $f$  (of the above form) can be expressed as a  $\mathbb{K}(x_1, \dots, x_s, n_1, \dots, n_r)$ -linear combination of the basis functions.

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Again, finitely many (modulo details) initial conditions suffice to fix  $f$ .

## Algebraic Setting

Write differential/difference equations in operator notation:

- ▶ shift operator  $S_v$ :  $S_v f(v) = f(v + 1)$
- ▶ partial derivative  $D_v$ :  $D_v f(v) = \frac{d}{dv} f(v)$
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$$(x^2 - 1)P_n''(x) + 2xP_n'(x) - n(n + 1)P_n(x) = 0$$

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**Example 2:** The three-term recurrence

$$nP_n(x) = (2n - 1)xP_{n-1}(x) - (n - 1)P_{n-2}(x)$$

translates to the operator

$$(n + 2)S_n^2 - (2n + 3)xS_n + (n + 1).$$

## Operator Algebra

Differential equations and recurrences are translated to skew polynomials.

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Even more general:

$$\partial_v a = \sigma(a) \partial_v + \delta(a)$$

where  $\sigma$  is an automorphism and  $\delta$  a  $\sigma$ -derivation, i.e.,

$$\delta(ab) = \sigma(a) \delta(b) + \delta(a) b.$$

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Such operators form an **Ore algebra**

$$\mathbb{K}(v, w, \dots) \langle \partial_v, \partial_w, \dots \rangle,$$

i.e., multivariate polynomials in the  $\partial$ 's with coefficients being rational functions in  $v, w, \dots$ , where  $\mathbb{K}$  is a computable field of characteristic 0 (i.e., containing  $\mathbb{Q}$ ).

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**Example:** The operators that we encountered with the Legendre polynomials live in the Ore algebra

$$\mathbb{K}(x, n) \langle D_x, S_n \rangle = \mathbb{K}(x, n) [D_x; 1, \frac{d}{dx}] [S_n; \sigma_n, 0].$$

## Annihilating Ideals

Let now  $\mathbb{D}$  be such an Ore algebra. The set

$$\text{Ann}_{\mathbb{D}} f := \{P \in \mathbb{D} : P \cdot f = 0\}$$

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are also valid equations for  $f$ . More generally,

$$\begin{aligned} P, Q \in \text{Ann}_{\mathbb{O}} f &\implies P + Q \in \text{Ann}_{\mathbb{O}} f \\ L \in \mathbb{O}, P \in \text{Ann}_{\mathbb{O}} f &\implies LP \in \text{Ann}_{\mathbb{O}} f \end{aligned}$$

which means that  $\text{Ann}_{\mathbb{O}} f$  is a **left ideal** in  $\mathbb{O}$ .

## Definition: $\partial$ -Finite Function

Let  $\mathbb{O} = \mathbb{K}(v, w, \dots) \langle \partial_v, \partial_w, \dots \rangle$  be an Ore algebra.

A function  $f(v, w, \dots)$  is  $\partial$ -finite w.r.t.  $\mathbb{O}$  if “all its shifts and derivatives”

$$\mathbb{O} \cdot f = \{P \cdot f : P \in \mathbb{O}\}$$

form a finite-dimensional  $\mathbb{K}(v, w, \dots)$ -vector space:

$$\dim_{\mathbb{K}(v, w, \dots)} (\mathbb{O} / \text{Ann}_{\mathbb{O}}(f)) < \infty.$$

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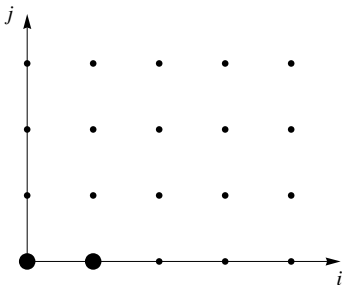
Let  $\mathbb{O} = \mathbb{K}(v, w, \dots) \langle \partial_v, \partial_w, \dots \rangle$  be an Ore algebra.

A function  $f(v, w, \dots)$  is  $\partial$ -finite w.r.t.  $\mathbb{O}$  if “all its shifts and derivatives”

$$\mathbb{O} \cdot f = \{P \cdot f : P \in \mathbb{O}\}$$

form a finite-dimensional  $\mathbb{K}(v, w, \dots)$ -vector space:

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We use (noncommutative) **Gröbner Bases** for certain operations on left ideals, e.g., for normal form computation modulo an ideal.

# Why $\partial$ -Finite Functions?

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3. These operations (closure properties) can be executed algorithmically.
4. Many elementary and special functions are covered.

## (Incomplete) List of $\partial$ -Finite Functions

ArcCsc, KelvinBei, HypergeometricPFQ, ExpIntegralE, ArcTanh, HankelH2, AngerJ, JacobiP, ChebyshevT, AiryBi, AiryAi, Sinc, Multinomial, CatalanNumber, QBinomial, CosIntegral, ArcSech, SphericalHankelH2, HermiteH, ExpIntegralEi, Beta, AiryBiPrime, SphericalBesselJ, Binomial, ParabolicCylinderD, Erfc, EllipticK, Fibonacci, QFactorial, Cos, Hypergeometric2F1, Erf, KelvinKer, HypergeometricPFQRegularized, Log, Factorial, BesselY, Cosh, CoshIntegral, ArcTan, ArcCoth, LegendreP, LaguerreL, EllipticE, SinhIntegral, Sinh, BetaRegularized, SphericalHankelH1, ArcSin, EllipticThetaPrime, Root, LucasL, AppellF1, FresnelC, LegendreQ, ChebyshevU, GammaRegularized, Erfi, HarmonicNumber, BesselJ, KelvinKei, ArithmeticGeometricMean, Exp, ArcCot, EllipticTheta, Hypergeometric0F1, EllipticPi, GegenbauerC, ArcCos, WeberE, FresnelS, EllipticF, ArcCosh, Subfactorial, QPochhammer, Gamma, StruveH, WhittakerM, ArcCsch, Hypergeometric1F1, SinIntegral, BesselJ, StruveL, ArcSec, Factorial2, KelvinBer, BesselK, ArcSinh, HankelH1, Sqrt, PolyGamma, HypergeometricU, AiryAiPrime, Sin,

That is nice, but we want (and can) do more...

What about integrals

$$\int_a^b f(x, \dots) dx$$

and sums

$$\sum_{n=a}^b f(n, \dots)$$

# Creative Telescoping

Method for doing integrals and sums  
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**Telescoping:** write  $f(n, k) = g(n, k + 1) - g(n, k)$ .

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**Creative Telescoping:** write

$$c_d(n)f(n + d, k) + \cdots + c_0(n)f(n, k) = g(n, k + 1) - g(n, k).$$

Summing from  $a$  to  $b$  yields a recurrence for  $F(n)$ :

$$c_d(n)F(n + d) + \cdots + c_0(n)F(n) = g(n, b + 1) - g(n, a).$$

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Method for doing integrals and sums  
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Consider the following integration problem:  $F(x) = \int_a^b f(x, y) \, dy$

**Telescoping:** write  $f(x, y) = \frac{d}{dy}g(x, y)$ .

Then  $F(x) = \int_a^b \left( \frac{d}{dy}g(x, y) \right) \, dy = g(x, b) - g(x, a)$ .

**Creative Telescoping:** write

$$c_d(x) \frac{d^d}{dx^d} f(x, y) + \cdots + c_0(x) f(x, y) = \frac{d}{dy}g(x, y).$$

Integrating from  $a$  to  $b$  yields a differential equation for  $F(x)$ :

$$c_d(x) \frac{d^d}{dx^d} F(x) + \cdots + c_0(x) F(x) = g(x, b) - g(x, a)$$

## Creative Telescoping, $\mathbb{O} = \mathbb{K}(n, k)\langle S_n, S_k \rangle$

$$\begin{aligned}c_d(n)f(n+d, k) + \cdots + c_0(n)f(n, k) &= g(n, k+1) - g(n, k) \\ &= (S_k - 1) \cdot g(n, k).\end{aligned}$$

Where should we look for a suitable  $g(n, k)$ ?

Note that there are trivial solutions like:

$$g(n, k) := \sum_{i=0}^{k-1} (c_d(n)f(n+d, i) + \cdots + c_0(n)f(n, i))$$

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A reasonable choice for where to look for  $g$  is  $\mathbb{O} \cdot f$ .

Then the task is to find  $P(n, S_n) = c_d(n)S_n^d + \cdots + c_0(n)$  and  $Q \in \mathbb{O}$  such that

$$(P - (S_k - 1)Q) \cdot f = 0 \quad \iff \quad P - (S_k - 1)Q \in \text{Ann}_{\mathbb{O}}(f).$$

## Creative Telescoping (Example 1)

Let  $F(n)$  denote the double sum over the trinomial coefficients

$$F(n) = \sum_{j=0}^n \sum_{i=0}^n \binom{n}{i, j, n-i-j} = \sum_{j=0}^n \sum_{i=0}^n \frac{n!}{i!j!(n-i-j)!}.$$

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Then the creative telescoping operator

$$CT = S_n - 3 + (S_i - 1) \frac{i}{n-i-j+1} + (S_j - 1) \frac{j}{n-i-j+1}$$

with  $CT \left( \binom{n}{i, j, n-i-j} \right) = 0$  implies that

$$F(n+1) = 3F(n).$$

## Creative Telescoping (Example 2)

The lattice Green's function of the square lattice is given by

$$P(z) = \int_0^1 \int_0^1 \frac{1}{(1 - xyz)\sqrt{1 - x^2}\sqrt{1 - y^2}} dx dy.$$

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The creative telescoping operator

$$(z^3 - z)D_z^2 + (3z^2 - 1)D_z + z + D_x \frac{y(1 - x^2)}{xyz - 1} + D_y \frac{yz(1 - y^2)}{xyz - 1}$$

that annihilates the integrand, certifies that  $P(z)$  satisfies the differential equation

$$(z^3 - z)P''(z) + (3z^2 - 1)P'(z) + zP(z) = 0.$$

## Creative Telescoping in Several Variables

In general, a creative telescoping operator has the form

$$P(\mathbf{v}, \partial_{\mathbf{v}}) + \Delta_1 Q_1(\mathbf{v}, \mathbf{w}, \partial_{\mathbf{v}}, \partial_{\mathbf{w}}) + \cdots + \Delta_m Q_m(\mathbf{v}, \mathbf{w}, \partial_{\mathbf{v}}, \partial_{\mathbf{w}})$$

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where  $\Delta_i = S_{w_i} - 1$  or  $\Delta_i = D_{w_i}$  (depending on the problem).

- ▶ Corresponds to an  $m$ -fold summation/integration problem.
- ▶  $\mathbf{w} = w_1, \dots, w_m$  are the summation/integration variables.
- ▶  $\mathbf{v} = v_1, v_2, \dots$  are the surviving parameters.
- ▶  $P(\mathbf{v}, \partial_{\mathbf{v}})$  is called the **telescoper**.
- ▶ The  $Q_i(\mathbf{v}, \mathbf{w}, \partial_{\mathbf{v}}, \partial_{\mathbf{w}})$  are called the **certificates**: they certify the correctness of the telescoper.

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Combine the two notions:

- ▶ Use  $\partial$ -finiteness for computations.
- ▶ Use holonomy for justifications.

## Holonomic Functions

Assume that  $f(x_1, \dots, x_s)$  depends only on continuous variables.  
Consider the **Weyl algebra**

$$\mathbb{W} = \mathbb{K}[x_1, \dots, x_s] \langle D_{x_1}, \dots, D_{x_s} \rangle.$$

Then  $f$  is holonomic if the left ideal  $\text{Ann}_{\mathbb{W}}(f)$  has dimension  $s$   
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→ This is why a creative telescoping operator always exists.

## $\partial$ -Finite and Holonomic Functions

**Theorem:** The function  $f(x_1, \dots, x_s)$  is holonomic if and only if it is  $\partial$ -finite.

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**Example:** The sequence  $\frac{1}{n^2+k^2}$  is  $\partial$ -finite but not holonomic.

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Eliminate all summation/integration variables from the annihilator ideal. Can be done, e.g., with Gröbner bases.

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- ▶ **Algorithms using Hermite reduction:**  
Allow to compute the telescoper without the certificates. So far, only available for special cases, e.g., hyperexponential functions.

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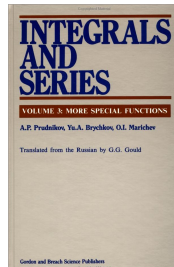
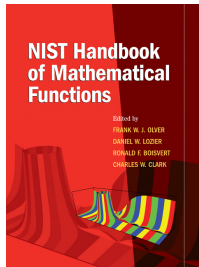
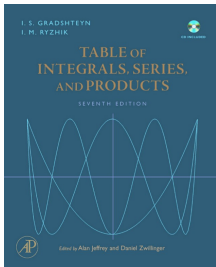
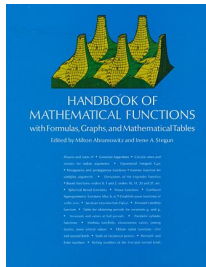
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3. Integrals and sums are treated by the method of creative telescoping.
4. The output is always given as an annihilating ideal, not as a closed form.

# Standard Application

# Special Function Identities



## Some Special Function Identities

$$\sum_{k=0}^n \binom{n}{k}^2 \binom{k+n}{k}^2 = \sum_{k=0}^n \binom{n}{k} \binom{k+n}{k} \sum_{j=0}^k \binom{k}{j}^3 \quad (1)$$

$$\int_0^{\infty} \frac{1}{(x^4 + 2ax^2 + 1)^{m+1}} dx = \frac{\pi P_m^{(m+\frac{1}{2}, -m-\frac{1}{2})}(a)}{2^{m+\frac{3}{2}}(a+1)^{m+\frac{1}{2}}} \quad (2)$$

$$e^{-x} x^{a/2} n! L_n^a(x) = \int_0^{\infty} e^{-t} t^{\frac{a}{2}+n} J_a(2\sqrt{tx}) dt \quad (3)$$

$$\int_{-\infty}^{\infty} \sum_{m=0}^{\infty} \sum_{n=0}^{\infty} \frac{H_m(x) H_n(x) r^m s^n e^{-x^2}}{m! n!} dx = \sqrt{\pi} e^{2rs} \quad (4)$$

$$\int_{-1}^1 (1-x^2)^{\nu-\frac{1}{2}} e^{iax} C_n^{(\nu)}(x) dx = \frac{\pi i^n \Gamma(n+2\nu) J_{n+\nu}(a)}{2^{\nu-1} a^{\nu} n! \Gamma(\nu)} \quad (5)$$

$$\frac{\sin(\sqrt{z^2 + 2tz})}{z} = \sum_{n=0}^{\infty} \frac{(-t)^n y_{n-1}(z)}{n!} \quad (6)$$

## Computer Proof of a Special Function Identity

$$e^{-x} x^{a/2} n! L_n^a(x) = \int_0^\infty e^{-t} t^{\frac{a}{2}+n} J_a(2\sqrt{tx}) dt.$$

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<< HolonomicFunctions.m
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```
CreativeTelescoping[Exp[-t]*t^(a/2+n)*BesselJ[a, 2*sqrt[t*x]]  
                  Der[t], {S[a], S[n], Der[x]}]
```

$$\{\{-2S_n + 2xD_x + (a + 2n + 2), \\ 4x^2D_x^2 + (4x^2 + 4x)D_x + (-a^2 + 2ax + 4nx + 4x), \\ 2xS_a^2 + (2ax + 2x^2 + 2x)D_x + (-a^2 + ax - a + 2nx + 2x)\}, \\ \{-2t, -4tx, -2tx\}\}$$

→ The annihilating ideals agree; check a few initial values.