

The size of the minimal automaton for an algebraic sequence

Eric Rowland

Hofstra University

Manon Stipulanti

University of Liège

Reem Yassawi

The Open University

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Algebraic sequences

\mathbb{F}_q denotes the finite field with q elements.

Let $s(n)_{n \geq 0}$ be a sequence of elements in \mathbb{F}_q .

$s(n)_{n \geq 0}$ is **algebraic** if there exists a nonzero polynomial $P(x, y) \in \mathbb{F}_q[x, y]$ such that $P(x, \sum_{n \geq 0} s(n)x^n) = 0$.

Combinatorial motivation: Integer sequences modulo p .

Example

Catalan numbers $C(n)_{n \geq 0} = 1, 1, 2, 5, 14, 42, 132, 429, \dots$

$F(x) = \sum_{n \geq 0} C(n)x^n$ satisfies $xy^2 - y + 1 = 0$ over \mathbb{Q} .

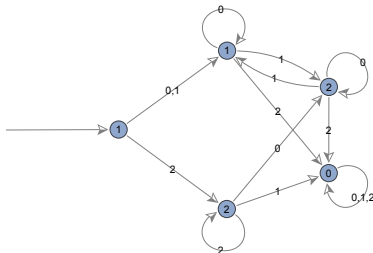
$F(x) = \sum_{n \geq 0} (C(n) \bmod 3)x^n$ satisfies $xy^2 + 2y + 1 = 0$ over \mathbb{F}_3 .

Automatic sequences

A sequence $s(n)_{n \geq 0}$ is **q -automatic** there is an automaton that outputs $s(n)$ when fed the base- q digits of n .

Convention in this talk: start with the least significant digit.

This automaton computes $C(n) \bmod 3$:



$C(9) = 4862 \equiv ? \pmod{3}$. Since $9 = 100_3$, $C(9) \equiv \boxed{2} \pmod{3}$.

$(C(n) \bmod 3)_{n \geq 0} = 1, 1, 2, 2, 2, 0, 0, 0, 2, 2, \dots$ is 3-automatic.

Theorem (Christol 1979/1980)

A sequence $s(n)_{n \geq 0}$ of elements in \mathbb{F}_q is algebraic if and only if it is q -automatic.

Two ways to represent such sequences: polynomials and automata.

How does the size of the automaton (number of states) depend on the x -degree (**height**) and y -degree (**degree**) of the polynomial?

Theorem (Bridy 2017)

Let $s(n)_{n \geq 0}$ be an algebraic sequence of elements in \mathbb{F}_q .
If its minimal polynomial has height h , degree d , and genus g , then the number of states in its minimal automaton is at most

$$(1 + o(1))q^{h+d+g-1},$$

where $o(1)$ tends to 0 as any of q, h, d, g gets large.

The genus satisfies $g \leq (h-1)(d-1)$; generically $g = (h-1)(d-1)$.

Corollary

The number of states is at most $(1 + o(1))q^{hd}$.

Can we get this bound without algebraic geometry? Yes.

Is the bound sharp? We suspect yes, but this is an open question.

How to construct an automaton?

Let $r \in \{0, 1, \dots, q-1\}$.

The **Cartier operator** Λ_r picks out every q th term, starting with $s(r)$:

$$\Lambda_r(s(n)_{n \geq 0}) := s(qn + r)_{n \geq 0}$$

Iteratively apply $\Lambda_0, \Lambda_1, \dots, \Lambda_{q-1}$ to $s(n)_{n \geq 0}$.

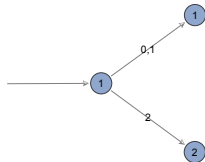
Create one state in the automaton for each distinct sequence.

Let $s(n) = (C(n) \bmod 3)$. $s(n)_{n \geq 0} = 1, 1, 2, 2, 2, 0, 0, 0, 2, \dots$

$\Lambda_0(s(n)_{n \geq 0}) = s(3n + 0)_{n \geq 0} = 1, 2, 0, 2, 1, 0, 0, 0, 0, \dots$ new!

$\Lambda_1(s(n)_{n \geq 0}) = s(3n + 1)_{n \geq 0} = 1, 2, 0, 2, 1, 0, 0, 0, 0, \dots = \Lambda_0(s(n)_{n \geq 0})$

$\Lambda_2(s(n)_{n \geq 0}) = s(3n + 2)_{n \geq 0} = 2, 0, 2, 1, 0, 0, 0, 0, 2, \dots$ new!



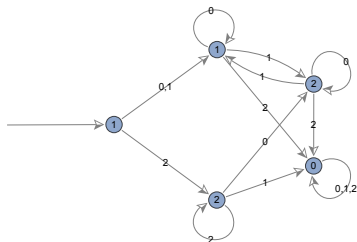
Label each state with the initial term of the corresponding sequence.

$$\Lambda_0(\Lambda_0(s(n)_{n \geq 0})) = 1, 2, 0, 2, 1, 0, 0, 0, 0, 2, \dots = \Lambda_0(s(n)_{n \geq 0})$$

$$\Lambda_1(\Lambda_0(s(n)_{n \geq 0})) = 2, 1, 0, 1, 2, 0, 0, 0, 0, 1, \dots \quad \text{new!}$$

$$\Lambda_2(\Lambda_0(s(n)_{n \geq 0})) = 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, \dots \quad \text{new!}$$

$$\Lambda_r(\Lambda_2(s(n)_{n \geq 0})) \quad \dots$$



A sequence is q -automatic if and only if this process terminates.

But we can't tell if sequences are equal from finitely many terms.

Use a different representation: diagonals of rational functions.

Theorem (Furstenberg 1967)

Let K be a field, and let $P(x, y) \in K[x, y]$ such that $\frac{\partial P}{\partial y}(0, 0) \neq 0$. If $F(x) \in K[[x]]$ satisfies $F(0) = 0$ and $P(x, F(x)) = 0$, then

$$F(x) = \mathcal{D} \left(\frac{y \frac{\partial P}{\partial y}(xy, y)}{P(xy, y)/y} \right).$$

The arguments xy arise from shearing the array of coefficients.

It will be more convenient to not shear. Then

$$F(x) = \mathcal{C} \left(\frac{y \frac{\partial P}{\partial y}(x, y)}{P(x, y)/y} \right)$$

where \mathcal{C} projects a Laurent series to the column $\langle x^i y^0 : i \geq 0 \rangle$.

Example

$\sum_{n \geq 0} (C(n) \bmod 3)x^n$ satisfies $xy^2 + 2y + 1 = 0$ over \mathbb{F}_3 .

$\sum_{n \geq 1} (C(n) \bmod 3)x^n$ is the y^0 column of

$$\begin{aligned} \frac{y \frac{\partial P}{\partial y}(x, y)}{P(x, y)/y} &= \frac{y(2xy + (2x + 2))}{(xy^2 + (2x + 2)y + x)/y} = 0x^0y^0 + 1x^0y^1 + 0x^0y^2 + 0x^0y^3 + 0x^0y^4 + 0x^0y^5 + \dots \\ &\quad + 0x^1y^{-1} + 1x^1y^0 + 0x^1y^1 + 2x^1y^2 + 0x^1y^3 + 0x^1y^4 + \dots \\ &\quad + 0x^2y^{-2} + 1x^2y^{-1} + 2x^2y^0 + 0x^2y^1 + 1x^2y^2 + 2x^2y^3 + \dots \\ &\quad + 0x^3y^{-3} + 1x^3y^{-2} + 1x^3y^{-1} + 2x^3y^0 + 0x^3y^1 + 1x^3y^2 + \dots \\ &\quad + 0x^4y^{-4} + 1x^4y^{-3} + 0x^4y^{-2} + 2x^4y^{-1} + 2x^4y^0 + 0x^4y^1 + \dots \\ &\quad + 0x^5y^{-5} + 1x^5y^{-4} + 2x^5y^{-3} + 0x^5y^{-2} + 0x^5y^{-1} + 0x^5y^0 + \dots \\ &\quad + \dots \end{aligned}$$

We have embedded $s(n)_{n \geq 0}$ into a bivariate series $\frac{S_0}{Q}$ where $Q = P/y$.
Can we compute $\Lambda_r(s(n)_{n \geq 0})$?

Define

$$\Lambda_r(x^n) = \begin{cases} x^{\frac{n-r}{q}} & \text{if } n \equiv r \pmod{q} \\ 0 & \text{otherwise} \end{cases}$$

and extend linearly to power series. Define $\Lambda_{r,s}$ analogously.

The map

$$\lambda_{r,0}(S) := \Lambda_{r,0}(SQ^{q-1})$$

on $\mathbb{F}_q[x, y]$ contains all information about $s \mapsto \Lambda_r(s)$ (and some extra):

$$\Lambda_r \mathcal{C} \left(\frac{S}{Q} \right) = \mathcal{C} \left(\frac{\Lambda_{r,0}(SQ^{q-1})}{Q} \right)$$

We construct an automaton by iterating $\lambda_{0,0}, \dots, \lambda_{q-1,0}$.

$(C(n) \bmod 3)_{n \geq 1}$ is a column of $\frac{S_0}{Q} := \frac{y(2xy+2x+2)}{(xy^2+(2x+2)y+x)/y}$.

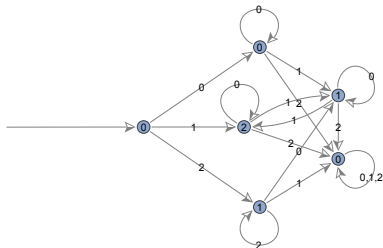
$$\lambda_{0,0}(S_0) = xy + x \quad \text{new!}$$

$$\lambda_{1,0}(S_0) = 2 \quad \text{new!}$$

$$\lambda_{2,0}(S_0) = y + 1 \quad \text{new!}$$

$$\lambda_{0,0}(xy + x) = xy + x = \lambda_{0,0}(S_0) \quad \dots$$

If two polynomials are equal, the corresponding sequences are equal.



The automaton may not be minimal.

Let $V := \langle x^i y^j : 0 \leq i \leq h \text{ and } 0 \leq j \leq d - 1 \rangle$. $\dim V = (h + 1)d$

Proposition

For each $r \in \{0, 1, \dots, q - 1\}$, we have $\lambda_{r,0}(S_0) \in V$.

For each $r \in \{1, \dots, q - 1\}$,

$$\lambda_{r,0}(V) \subseteq \langle x^i y^j : 0 \leq i \leq h - 1 \text{ and } 0 \leq j \leq d - 1 \rangle$$

which has dimension hd .

Corollary:

The constructed automaton has at most $q^{hd} + |\text{orb}_{\lambda_{0,0}}(S_0)|$ states.

It remains to bound $|\text{orb}_{\lambda_{0,0}}(S_0)|$.

Certain orders of the basis for V show that $\lambda_{0,0}$ is highly structured.

Example

Let $q = 3$, $h = 2$, $d = 4$, and

$$P = (x^2 + x + 2)y^4 + xy^3 + (2x + 1)y^2 + (x^2 + 1)y + 2x^2 + x.$$

Basis:

$$(x^1y^0, x^1y^1, x^1y^2, x^0y^1, x^0y^2, x^0y^0, x^1y^3, x^0y^3, x^2y^0, x^2y^1, x^2y^2, x^2y^3).$$

Matrix for $\lambda_{0,0}$:

$$\begin{bmatrix} 1 & 1 & 1 & 2 & 1 & 2 & 0 & 0 & 2 & 2 & 0 & 0 \\ 1 & 2 & 1 & 2 & 2 & 2 & 1 & 2 & 1 & 1 & 1 & 2 \\ 2 & 2 & 1 & 2 & 1 & 2 & 1 & 2 & 1 & 1 & 1 & 1 \\ & & & 1 & 2 & 1 & & 1 & & & & \\ & & & 0 & 1 & 1 & & 1 & & & & \\ & & & & & & 1 & & & & & \\ & & & & & & & 2 & 2 & & & 1 \\ & & & & & & & & 1 & & & \\ & & & & & & & & & 1 & 1 & 0 \\ & & & & & & & & & 2 & 1 & 0 & 1 \\ & & & & & & & & & 1 & 0 & 0 & 2 \\ & & & & & & & & & & & & 1 \end{bmatrix}$$

Basis of V :

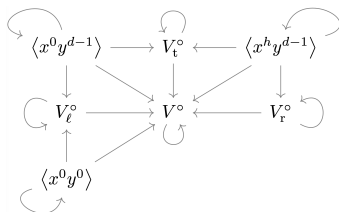
$x^0 y^{d-1}$	$x^1 y^{d-1}$...	$x^{h-1} y^{d-1}$	$x^h y^{d-1}$

$x^0 y^{d-2}$	$x^1 y^{d-2}$...	$x^{h-1} y^{d-2}$	$x^h y^{d-2}$
⋮				
$x^0 y^1$	$x^1 y^1$...	$x^{h-1} y^1$	$x^h y^1$

$x^0 y^0$	$x^1 y^0$...	$x^{h-1} y^0$	$x^h y^0$

Theorem

Under applications of $\lambda_{0,0}$ on V , information flows as follows.



The left, right, and top subspaces are affected only by themselves. Since $|V^\circ| = q^{(h-1)(d-1)}$, we show the borders contribute $\leq q^{h+d-1}$.

The left, right, and top subspaces are essentially univariate.

Fix $R \in \mathbb{F}_q[z]$. How big are orbits under $\lambda_0(S) := \Lambda_0(SR^{q-1})$?
This is “just” a linear transformation.

Example

Let $q = 3$ and $R = (z^2 + 1)(z^3 + z^2 + 2) \in \mathbb{F}_3[z]$.

Compute $\text{orb}_{\lambda_0}(S)$ from each $S \in \mathbb{F}_3[z]$ with $\deg S \leq \deg R$.

Period lengths that occur: $\{1, 2, 3, 6\}$

Example

Let $q = 3$ and $R = (z^2 + 1)(z^4 + z + 2) \in \mathbb{F}_3[z]$.

Period lengths: $\{1, 2, 4\}$

Consider all polynomials R with fixed degree.

Surprising fact: The maximal period length doesn't depend on q .

Theorem

Let $R \in \mathbb{F}_q[z]$ such that $R \neq 0$, $z \nmid R$, and R is square-free. Let $cR_1 \cdots R_m$ be its factorization into irreducibles, and let

$$\ell = \text{lcm}(\deg R_1, \dots, \deg R_m).$$

Then $\lambda_0^\ell(S) = S$ for all $S \in \mathbb{F}_q[z]$ with $\deg S \leq \deg R$.

The upper bound is achieved when ℓ is maximized, subject to $\deg R_1 + \cdots + \deg R_m = \deg R$.

The **Landau function** $L(n)$ is the maximum value of $\text{lcm}(n_1, \dots, n_m)$ over all integer partitions (n_1, \dots, n_m) of n . Also arises in Bridy's proof.

Corollary

The number of states is at most

$$q^{hd} + q^{(h-1)(d-1)} L(h)L(d)^2 + \left\lceil \log_q \max(h, d, q) \right\rceil.$$

Asymptotically...

Landau (1903): $\log L(n) \sim \sqrt{n \log n}$

Massias–Nicolas–Robin (1988): $L(n) \leq e^{(1+o(1))\sqrt{n \log n}}$

Corollary

The number of states is at most $(1 + o(1))q^{hd}$.







Example

The factor $1 + o(1)$ cannot be removed. Let $q = 2$ and

$$P = (x^3 + x^2 + 1)y^3 + (x^3 + 1)y^2 + (x^3 + x^2 + x + 1)y + x^3 + x^2$$

with $h = 3$ and $d = 3$. The number of states is $532 > 512 = q^{hd}$.

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